



Paraconsistent Reasoning for OWL 2



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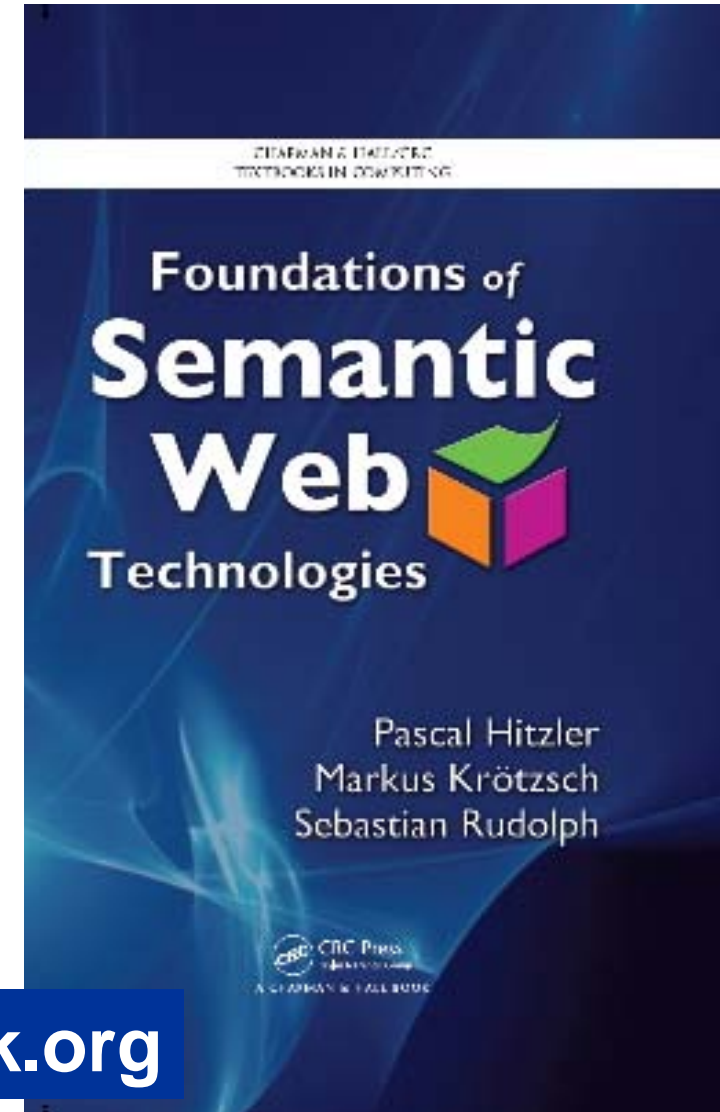
Kno.e.sis Center
Wright State University, Dayton, OH
<http://www.knoesis.org/pascal/>

Pascal Hitzler, Markus Krötzsch,
Sebastian Rudolph

Foundations of Semantic Web
Technologies
Chapman & Hall/CRC, 2009

Grab a flyer!

<http://www.semantic-web-book.org>





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- Paraconsistent semantics which can deal with inconsistent OWL 2 DL ontologies, based on 4-valued logic
- PolyTime transformation into OWL 2 DL (i.e. standard OWL reasoner can be used)
- Tractability of OWL 2 EL and OWL 2 QL retained under paraconsistent semantics
- Prototype implementations exist

- **Motivation**
- The Semantics
- PolyTime Transformation to classical OWL
- Tractable Fragments
- Implementation
- What needs to be done

Idea:

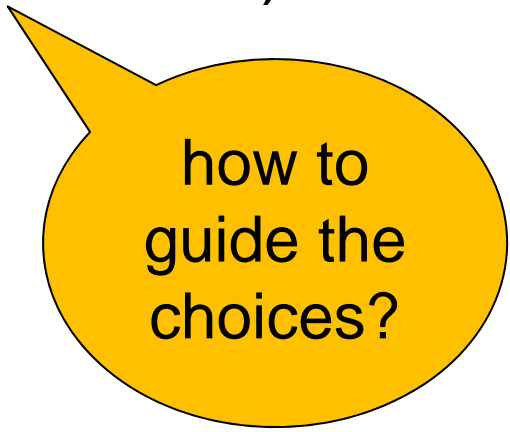
Rather than repairing inconsistencies
change semantics to provide for them

Rationales:

- Inconsistencies occur naturally, i.e. they do not need repair but dealing with
- Repair is too difficult or too expensive

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- Use 4 truth values instead of 2:
 - true
 - false
 - undefined / unknown
 - overdefined / both / inconsistent
- Make reasonable truth tables (choices exist)
- Lift semantics to the 4 truth values



how to
guide the
choices?

- Does not increase OWL 2 computational complexity
- Tractable fragments of OWL 2 remain tractable
- Polytime transformations to 2-valued semantics
- Standard equivalences from 2-valued semantics still hold (e.g. DeMorgan)

Constructor Syntax	Semantics
A	$A^I = \langle P, N \rangle$, where $P, N \subseteq \Delta^I$
R	$R^I \subseteq \Delta^I \times \Delta^I$
o	$o^I \in \Delta^I$
\top	$\langle \Delta^I, \emptyset \rangle$
\perp	$\langle \emptyset, \Delta^I \rangle$
$C_1 \sqcap C_2$	$\langle P_1 \cap P_2, N_1 \cup N_2 \rangle$, if $C_i^I = \langle P_i, N_i \rangle$ for $i = 1, 2$
$C_1 \sqcup C_2$	$\langle P_1 \cup P_2, N_1 \cap N_2 \rangle$, if $C_i^I = \langle P_i, N_i \rangle$ for $i = 1, 2$
$\neg C$	$(\neg C)^I = \langle N, P \rangle$, if $C^I = \langle P, N \rangle$
$\exists R.C$	$\langle \{x \mid \exists y, (x, y) \in R^I \text{ and } y \in proj^+(C^I)\}, \{x \mid \forall y, (x, y) \in R^I \text{ implies } y \in proj^-(C^I)\} \rangle$
$\forall R.C$	$\langle \{x \mid \forall y, (x, y) \in R^I \text{ implies } y \in proj^+(C^I)\}, \{x \mid \exists y, (x, y) \in R^I \text{ and } y \in proj^-(C^I)\} \rangle$

Constructor	Semantics
$\geq nR.C$	$\langle \{x \mid \#(y.(x, y) \in R^I \wedge y \in proj^+(C^I)) \geq n\}, \{x \mid \#(y.(x, y) \in R^I \wedge y \notin proj^-(C^I)) < n\} \rangle$
$\leq nR.C$	$\langle \{x \mid \#(y.(x, y) \in R^I \wedge y \notin proj^-(C^I)) \leq n\}, \{x \mid \#(y.(x, y) \in R^I \wedge y \in proj^+(C^I)) > n\} \rangle$
$\{o_1, \dots, o_n\}$	$\langle \{o_1^I, \dots, o_n^I\}, N \rangle$, where $N \subseteq \Delta^I$

- Nominals are basically treated like named classes; just that their positive part is fixed.

Axiom Name	Syntax	Semantics
material inclusion	$C_1 \mapsto C_2$	$\Delta^I \setminus proj^-(C_1^I) \subseteq proj^+(C_2^I)$
internal inclusion	$C_1 \sqsubset C_2$	$proj^+(C_1^I) \subseteq proj^+(C_2^I)$
strong inclusion	$C_1 \rightarrow C_2$	$proj^+(C_1^I) \subseteq proj^+(C_2^I)$ and $proj^-(C_2^I) \subseteq proj^-(C_1^I)$
individual assertions	$C(a)$ $R(a, b)$	$a^I \in proj^+(C^I)$ $(a^I, b^I) \in R^I$

- Three choices for resolving class inclusion exist – they differ in semantic strength.
- Note roles are 2-valued!

$\varphi \mapsto \psi$ is definable as $\neg\varphi \vee \psi$. (Material Implication)

$\varphi \supset \psi$ evaluates to
$$\begin{cases} \psi & \text{if } \varphi \in \{t, \ddot{\top}\} \\ t & \text{if } \varphi \in \{f, \ddot{\perp}\} \end{cases}$$
 (Internal Implication)

$\varphi \rightarrow \psi$ is definable as $(\varphi \supset \psi) \wedge (\neg\psi \supset \neg\varphi)$ (Strong Implication)

- Material is the weakest of the three
- Internal satisfies the deduction theorem
- Strong additionally satisfies contraposition
- When to choose which implication is still unclear

- Every knowledge base has a four-valued model
- (under any mix of class inclusions)
- Caveat: Have to remove \top and \perp first

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- (in fact, linear)
- Standard method:
Rewrite $\neg A$ to new class name A'
and treat them separately.
Then lift this transformation over the structure of
formulae $\Rightarrow \pi$.
- $KB \models_4 \alpha$ iff $\pi(KB) \models_2 \pi(\alpha)$

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- OWL 2 EL: $\pi(K)$ is in this fragment if
 - K is and
 - only internal inclusion is used
- Horn-SHIQ: Transformation needs modification.
 $\pi_H(K)$ is in Horn-SHIQ if
 - K is and
 - only internal inclusion is used
- OWL 2 QL: Can be retained, but semantics needs modifications

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- By Steffen Stadtmüller
Universität Karlsruhe



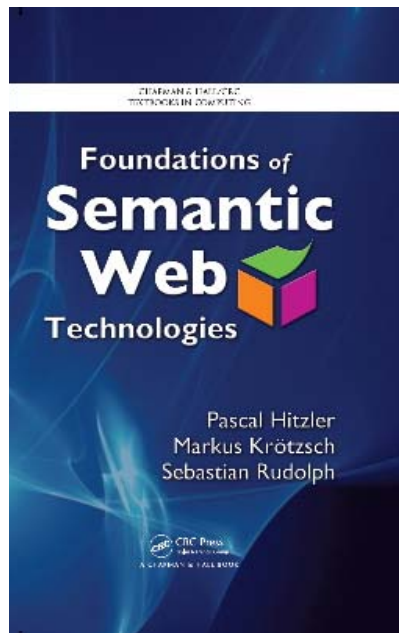
Part of the RaDON plugin of the NeOn Toolkit
See ESWC2009 Demo paper on RaDON
<http://www.neon-toolkit.org>

- By Fred Maier, Florida Institute for Human and
Machine Cognition

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- Extensive testing of the tool
 - efficiency
 - comparison with other approaches (e.g. debugging)
- Investigate which class inclusion to use in which context
 - automatic classification?
 - iterative weakening?
 - are there other ways to resolve this?
- Apply in realistic application scenarios

Thanks!



In a Nutshell:

- Paraconsistent semantics for OWL 2 DL based on 4-valued logic
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